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Quantum Secret Sharing

We provide a secure quantum carrier for distributing a secret (classical symbol encoded into a state or a quantum state) among n parties according to a (k, n) threshold scheme, where $2k - 1$ is a prime number.

- Re-usable quantum carrier
- Make use of the results of every round
- Predetermined message can be sent securely

Quantum Carrier: 

(2,3) Scheme, Basic properties

Using this encoding each two participants can retrieve the secret:

$$\begin{aligned} 0 &\rightarrow |\bar{0}\rangle = |000\rangle + |111\rangle + |222\rangle \\ 1 &\rightarrow |\bar{1}\rangle = |012\rangle + |120\rangle + |201\rangle \\ 2 &\rightarrow |\bar{2}\rangle = |021\rangle + |102\rangle + |210\rangle \end{aligned}$$

Using CNOT gates

$$C_{a,1}^0 C_{a,2}^1 C_{a,3}^2 |i\rangle_a |\bar{s}\rangle_{1,2,3} = |i\rangle_a |\bar{i+s}\rangle_{1,2,3}$$

$$C_{b,1}^1 C_{b,2}^2 C_{b,3}^3 |\bar{i}\rangle_{a,b,c} |\bar{s}\rangle_{1,2,3} = |\bar{i}\rangle_{b_1,b_2,b_3} |\bar{i+s}\rangle_{1,2,3}$$

Performing Hadamard gate

$$(H|j\rangle = \sum_k \xi^{jk} |k\rangle \text{ with } \xi = e^{\frac{2\pi i}{3}}$$

We have:

$$H^{\otimes 3} |\bar{s}\rangle = \sum_x \xi^{-sx} |\bar{x}\rangle$$

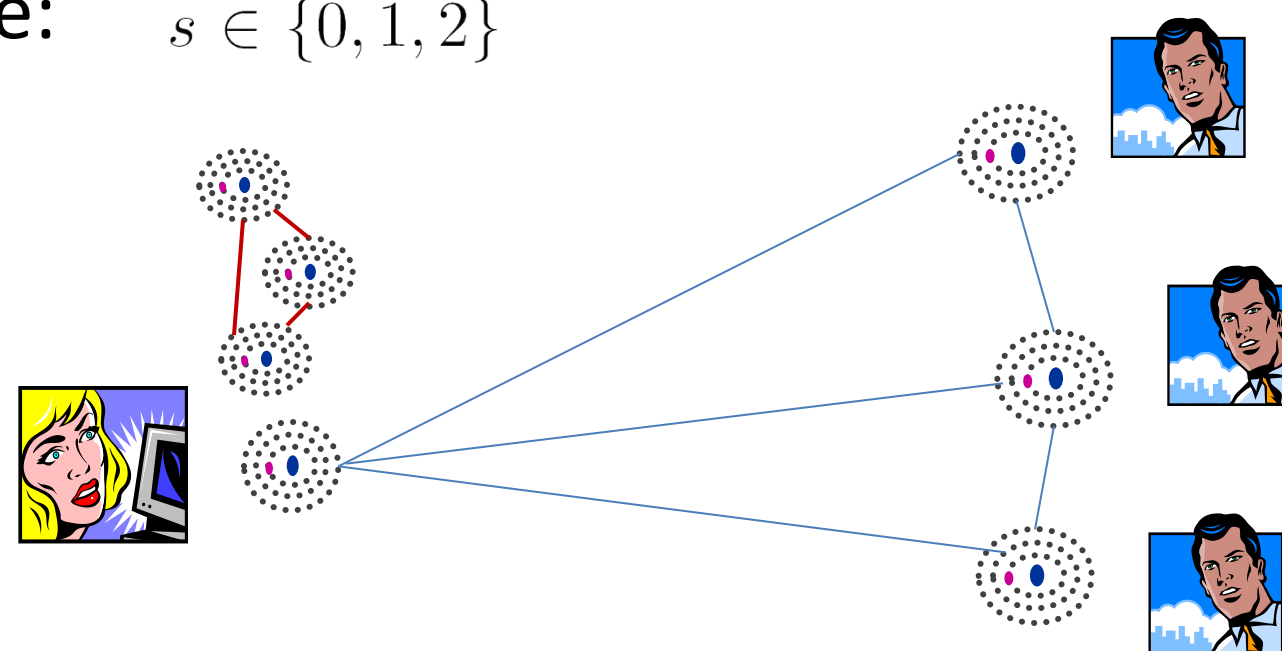
So, on our Carrier:

$$|\Phi\rangle = \sum_i |i\rangle_A |\bar{i}\rangle_{B_1,B_2,B_3}$$

$$H^{\otimes 4} |\Phi\rangle = |\Phi\rangle$$

Scheme

I) Alice chooses the classical secret and prepares the quantum code: $s \in \{0, 1, 2\}$



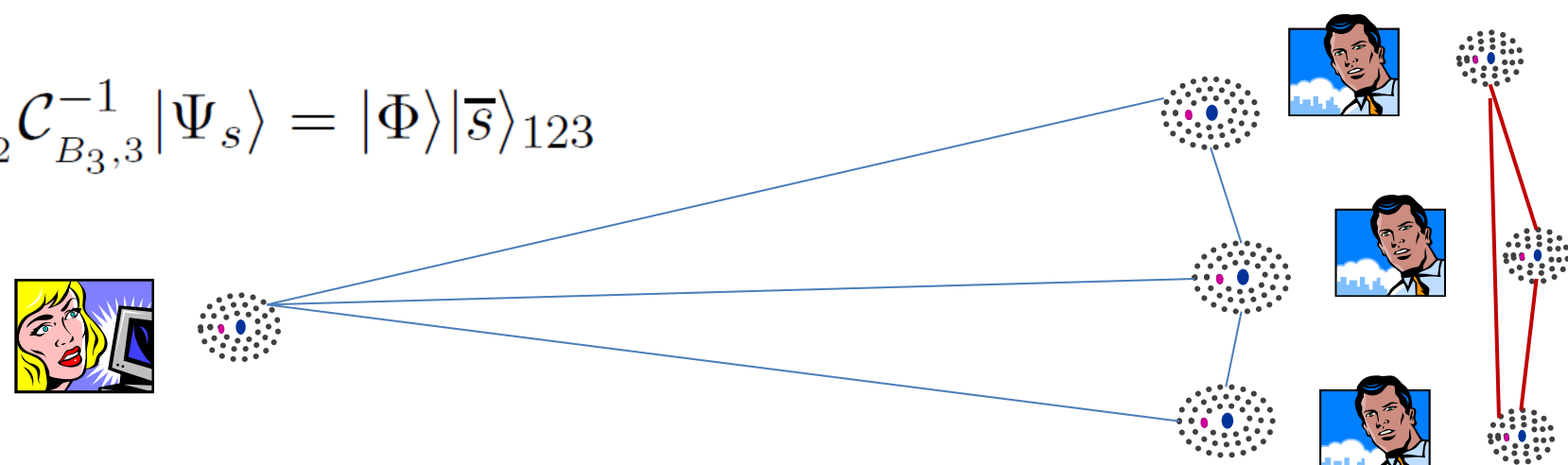
II) Alice hooks the quantum code to the quantum carrier:

$$|\Psi_s\rangle = C_{A,1}^0 C_{A,2}^1 C_{A,3}^2 |\Phi\rangle |\bar{s}\rangle_{123} = \sum_{i=0}^2 |i\rangle_A |\bar{i}\rangle_{B_1 B_2 B_3} |\bar{i+s}\rangle_{123}$$

And sends the qutrits to the participants.

III) Each participant locally unhooks his qutrit from quantum carrier:

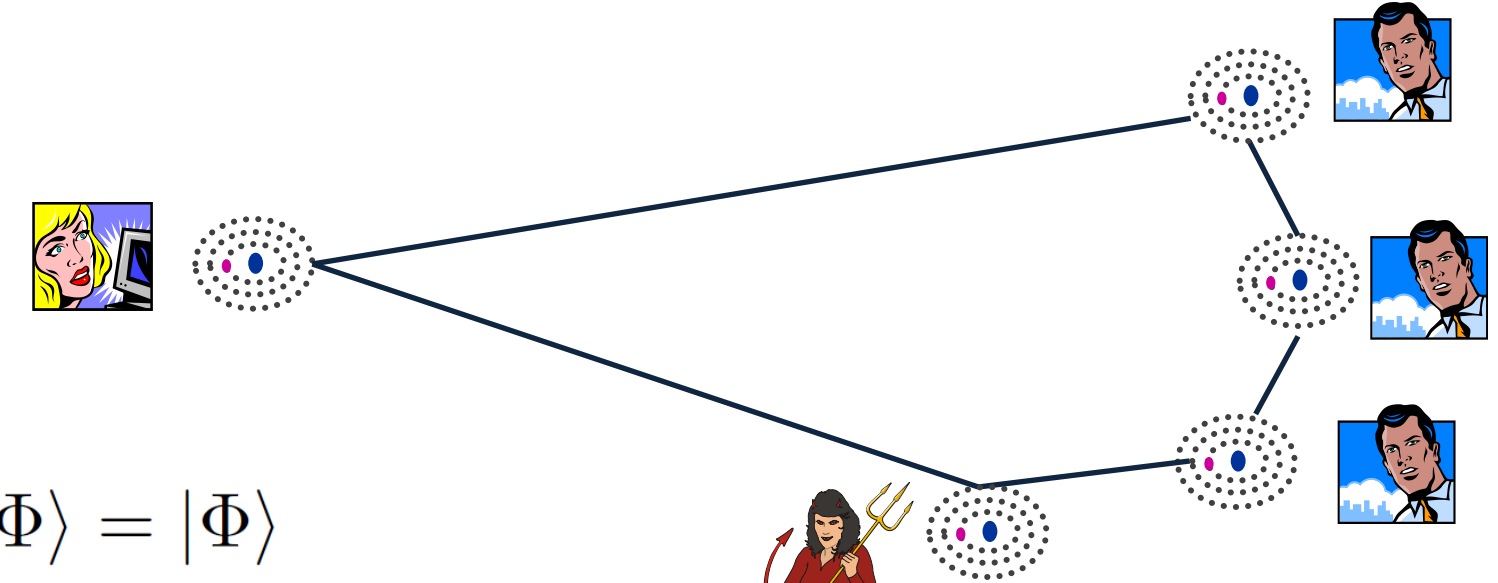
$$C_{B_1,1}^{-1} C_{B_2,2}^{-1} C_{B_3,3}^{-1} |\Psi_s\rangle = |\Phi\rangle |\bar{s}\rangle_{123}$$



IV) Participants measure locally. With collaboration of each two of them the secret can be retrieved.

V) Alice and participants perform Hadamard gates.

Security



$$H^{\otimes 4} |\Phi\rangle = |\Phi\rangle$$

$$|\Phi'\rangle = (H^{\otimes 4} \otimes U) \sum_j |j, \bar{j}\rangle |\phi_j\rangle$$

Betray of a participant:

$$\begin{aligned} \sum_{i,j} |i, j + i, j + 2i, j\rangle_{A,B_1,B_2,B_3} |\phi_j\rangle_{B'_3} \\ \rho_{B_3,B'_3,1,2,3} = \sum_j |j\rangle\langle j|_{B_3} \otimes |\phi_j\rangle\langle\phi_j|_{B'_3} \otimes \sum_i |\bar{i+s}\rangle\langle\bar{i+s}| \end{aligned}$$

Scheme

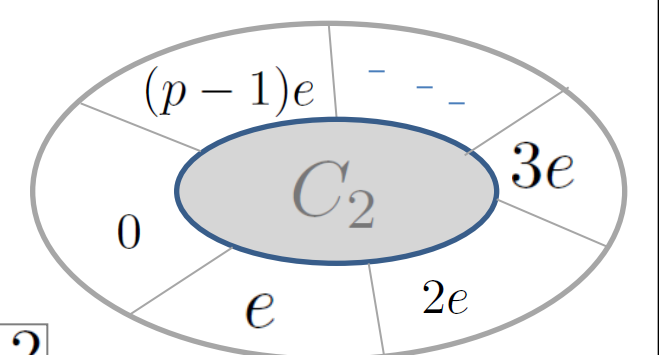
Using CSS codes:

$$C_2 \subset C_1$$

$$k_1 - k_2 = 1$$

$$|\bar{w}\rangle = \frac{1}{\sqrt{p^{k_2}}} \sum_{c_2 \in C_2} |c_2 + w\rangle \quad w \in C_1 \quad w = v + ie$$

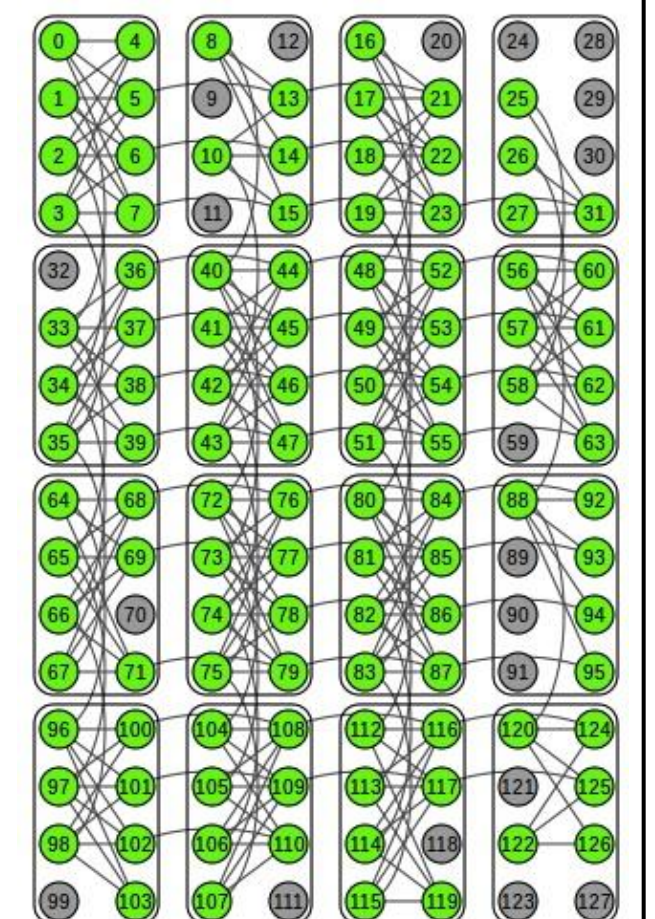
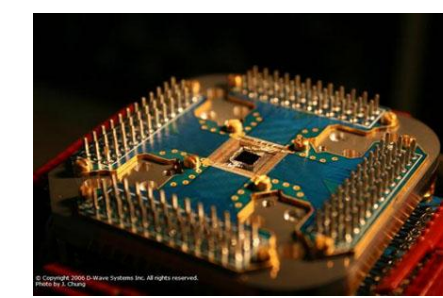
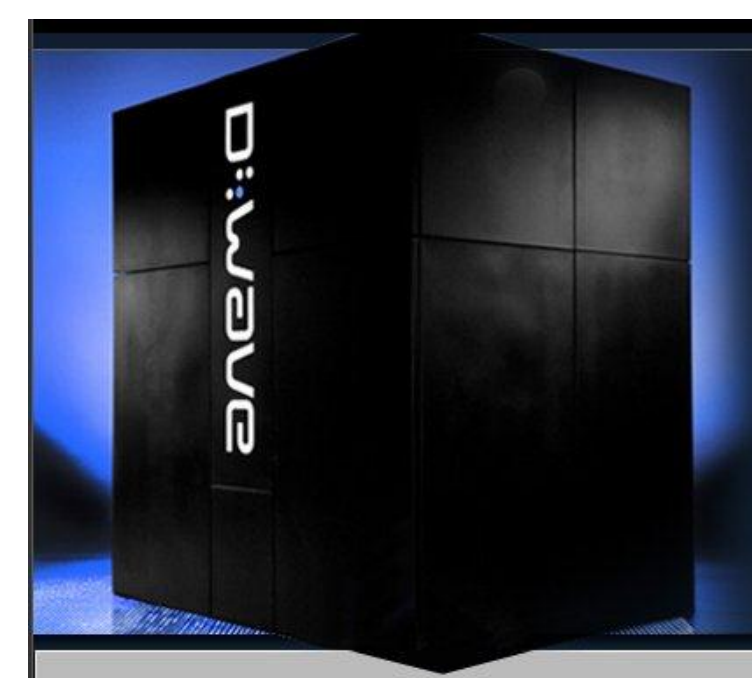
$$|\bar{i}\rangle = \sum_{v \in C_2} |v + ie\rangle$$



$$\begin{aligned} e_i \cdot e_j &= 0 & 1 \leq i, j \leq k-2 \\ e \cdot e_j &= 0 & 0 \leq j \leq k-2 \\ e \cdot e &= -1 \end{aligned}$$

Ref: <http://arxiv.org/abs/1007.2778>, In Press(IJQI)

D-Wave one Benchmarking



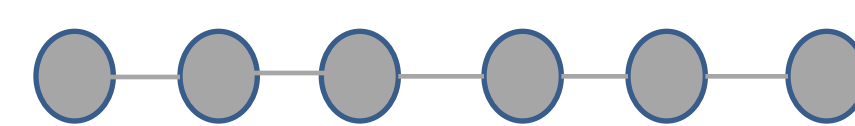
Quantum Adiabatic evolution:

$$H(t) = A(t) \sum \sigma_x + B(t) H_{Ising \text{ model}}$$

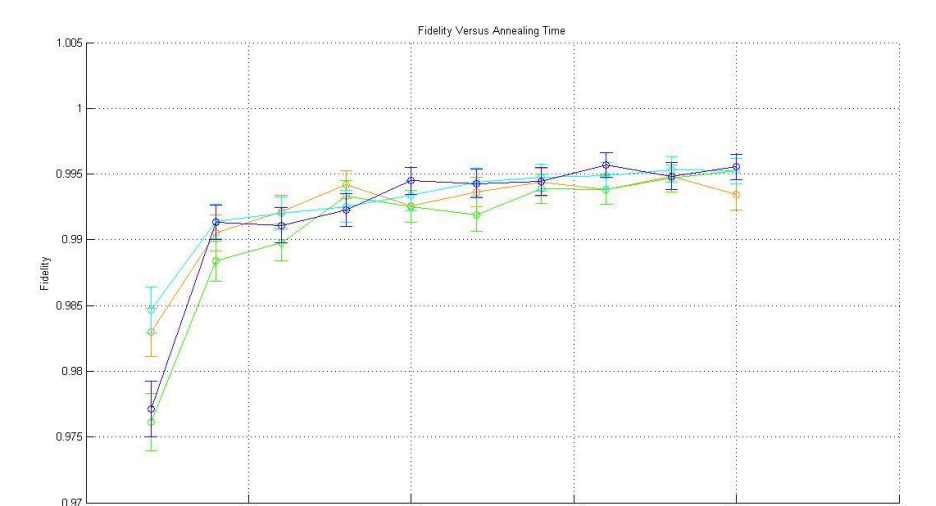
Input & output: h & J Matrices $S_i \in \{-1, 1\}$

$$\text{finding } S^*: \text{ argmin}_s \left\{ \sum h_i S_i + \sum J_{ij} S_i S_j \right\}$$

1D Ising model



$$H = -J \sum g \hat{\sigma}_i^x + \hat{\sigma}_i^z \hat{\sigma}_i^z$$



$$\text{Exact spectrum } \varepsilon_k = 2J \sqrt{(1 + g^2 - 2g \cos(k))}$$

$$k \propto \frac{\pi}{N}$$